expkv

an expandable (key)=(value) implementation

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Abstract

expkv provides a small interface for $\langle key \rangle = \langle value \rangle$ parsing. The parsing macro is *fully expandable*, the $\langle code \rangle$ of your keys might be not. expkv is *swift*, close to the fastest $\langle key \rangle = \langle value \rangle$ implementation. However it is the fastest which copes with active commas and equal signs and doesn't strip braces accidentally.

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1 Documentation

expkv provides an expandable $\langle key \rangle = \langle value \rangle$ parser. The $\langle key \rangle = \langle value \rangle$ pairs should be given as a comma separated list and the separator between a $\langle key \rangle$ and the associated $\langle value \rangle$ should be an equal sign. Both, the commas and the equal signs, might be of category 12 (other) or 13 (active). To support this is necessary as for example babel turns characters active for some languages, for instance the equal sign is turned active for Turkish.

expkv is usable as generic code, as a LATEX package or as a ConTEXt module. To use it, just use one of:

```
\input expkv % plainTeX
\usepackage{expkv} % LaTeX
\usemodule[expkv] % ConTeXt
```

Both the LATEX package and the ConTeXt module don't do more than expkv.tex, except calling \ProvidesPackage and setting things up such that expkv.tex will use \ProvidesFile, or printing some status information. The ConTeXt support is not thoroughly tested, though (since I don't use ConTeXt myself I don't know if there are additional pitfalls I wasn't aware of).

In the expkv family are other packages contained which provide additional functionality. Those packages currently are:

```
expkvider a key-defining frontend for expkv using a \langle key \rangle = \langle value \rangle syntax expkvides define expandable \langle key \rangle = \langle value \rangle macros using expkv expkviopt parse package and class options with expkv
```

Note that while the package names are stylised with a vertical rule, their names are all lower case with a hyphen (*e.g.*, expkv-def).

A list of concise comparisons to other $\langle key \rangle = \langle value \rangle$ packages is contained in subsection 1.7.

1.1 Setting up Keys

expkv provides a rather simple approach to setting up keys, similar to keyval. However there is an auxiliary package named expkvider which provides a more sophisticated interface, similar to well established packages like pgfkeys or l3keys.

Keys in $\exp_{\mathbf{k}\mathbf{v}}$ (as in almost all other $\langle key \rangle = \langle value \rangle$ implementations) belong to a *set* such that different sets can contain keys of the same name. Unlike many other implementations $\exp_{\mathbf{k}\mathbf{v}}$ doesn't provide means to set a default value, instead we have keys that take values and keys that don't (the latter are called NoVal keys by $\exp_{\mathbf{k}\mathbf{v}}$), but both can have the same name (on the user level).

The following macros are available to define new keys. Those macros containing "def" in their name can be prefixed by anything allowed to prefix \def (but don't use \outer, keys defined with it won't ever be usable), prefixes allowed for \let can prefix those with "let" in their name, accordingly. Neither \left\(set\right\) nor \left\(key\right\) are allowed to be empty for new keys. \left\(set\right\) will be used as is inside of \csname \ldot\.\end{csname} and \left\(key\right\) will get \detokenized. Also \left\(set\right\) should not contain an explicit \par token.

\ekvdef

 $\left(\left(set \right) \right) \left(\left(code \right) \right)$

Defines a $\langle key \rangle$ taking a value in a $\langle set \rangle$ to expand to $\langle code \rangle$. In $\langle code \rangle$ you can use #1 to refer to the given value.

Example: Define text in foo to store the value inside \foo@text: \protected \long \ekvdef \{ foo \} \{ text \} \\ \def \foo@text \{\#1\} \}

\ekvdefNoVal

 $\verb|\ekvdefNoVal{$\langle set \rangle$} {\langle key \rangle} {\langle code \rangle}$

Defines a no value taking $\langle key \rangle$ in a $\langle set \rangle$ to expand to $\langle code \rangle$.

Example: Define bool in foo to set \iffoo@bool to true:
\protected\ekvdefNoVal{foo}{bool}{\foo@booltrue}

\ekvlet

 $\verb|\ekvlet{\langle set \rangle}| \langle key \rangle| \langle cs \rangle|$

Let the value taking $\langle key \rangle$ in $\langle set \rangle$ to $\langle cs \rangle$, there are no checks on $\langle cs \rangle$ enforced, but the code should expect the value as a single braced argument directly following it.

Example: Let cmd in foo do the same as \foo@cmd:

\ekvlet { foo } {cmd} \foo@cmd

\ekvletNoVal

 $\verb|\ekvletNoVal{|} (set)| {\langle key \rangle} {\langle cs \rangle}$

Let the no value taking $\langle key \rangle$ in $\langle set \rangle$ to $\langle cs \rangle$, it is not checked whether $\langle cs \rangle$ exists or that it takes no parameter.

Example: See above.

\ekvletkv

 $\verb|\ekvletkv|{\langle set \rangle}|{\langle key \rangle}|{\langle set 2 \rangle}|{\langle key 2 \rangle}|$

Let the $\langle key \rangle$ in $\langle set \rangle$ to $\langle key2 \rangle$ in $\langle set2 \rangle$, it is not checked whether that second key exists (but take a look at \ekvifdefined).

Example: Let B in bar be an alias for A in foo:

ekvletkv{bar}{B}{foo}{A}

\ekvletkvNoVal

 $\verb|\ekvletkvNoVal{$\langle set \rangle$} {\langle key \rangle} {\langle set 2 \rangle} {\langle key 2 \rangle}$

Let the $\langle key \rangle$ in $\langle set \rangle$ to $\langle key2 \rangle$ in $\langle set2 \rangle$, it is not checked whether that second key exists (but take a look at \ekvifdefinedNoVal).

Example: See above.

\ekvdefunknown

 $\verb|\ekvdefunknown{$\langle set \rangle$} {\langle code \rangle$}$

By default an error will be thrown if an unknown $\langle key \rangle$ is encountered. With this macro you can define $\langle code \rangle$ that will be executed for a given $\langle set \rangle$ when an unknown $\langle key \rangle$ with a $\langle value \rangle$ was encountered instead of throwing an error. You can refer to the given $\langle value \rangle$ with #1 and to the unknown $\langle key \rangle$'s name with #2 in $\langle code \rangle$.\(^1\) \ekvdefunknown and \ekvredirectunknown are mutually exclusive, you can't use both.

Example: Also search bar for undefined keys of set foo: \long\ekvdefunknown\{foo}\{\ekvset\{bar}\{\pmu2=\{\pmu1}\}\}

tiong tenunejunknown () oo j t tenuser tour j (#2=(#1 j

¹That order is correct, this way the code is faster.

This example differs from using \ekvredirectunknown{foo}{bar} (see below) in that also the unknown-key handler of the bar set will be triggered, error messages for undefined keys will look different, and this is slower than using \ekvredirectunknown.

\ekvdefunknownNoVal

 $\verb|\ekvdefunknownNoVal{$\langle set \rangle$} {\langle code \rangle$}$

As already explained for \ekvdefunknown, expkv would throw an error when encountering an unknown \langle key \rangle. With this you can instead let it execute \langle code \rangle if an unknown NoVal \langle key \rangle was encountered. You can refer to the given \langle key \rangle with #1 in \langle code \rangle. \langle ekvdefunknownNoVal and \langle ekvredirectunknownNoVal are mutually exclusive, you can't use both.

Example: Also search bar for undefined keys of set foo: \\ekvdefunknownNoVal\{foo}\{\ekvset\{bar\}\{\pm1\}\\\}

\ekvredirectunknown

 $\ensuremath{\ensuremath{\mbox{\sc hown}}} {\ensuremath{\mbox{\sc hown}}}} {\ensuremath{\mbox{\sc hown}}} {\ensuremath{\mbox{\sc hown}}}} {\ensuremath{\mbox{\sc hown}}} {\ensuremath{\mb$

This is a short cut to set up a special \ekvdefunknown for $\langle set \rangle$ that will check each set in the comma separated $\langle set-list \rangle$ for the unknown $\langle key \rangle$. You can't use prefixes (so no \long or \protected) with this macro, the resulting unknown-key handler will always be \long. The first set in the $\langle set-list \rangle$ has highest priority. Once the $\langle key \rangle$ is found the remaining sets are discarded, if the $\langle key \rangle$ isn't found in any set an error will be thrown eventually. Note that the error messages are affected by the use of this macro, in particular, it isn't checked whether a NoVal key of the same name is defined in order to throw an unwanted value error. \ekvdefunknown and \ekvredirectunknown are mutually exclusive, you can't use both.

Example: For every key not defined in the set foo also search the sets bar and baz: $\ensuremath{\mbox{\it ekvredirectunknown}\{foo\}\{bar,\ baz\}$

\ekvredirectunknownNoVal

\ekvredirectunknownNoVal $\{\langle set \rangle\}\{\langle set-list \rangle\}$

This behaves just like \ekvredirectunknown and does the same but for the NoVal keys. Again no prefixes are supported. Note that the error messages are affected by the use of this macro, in particular, it isn't checked whether a normal key of the same name is defined in order to throw a missing value error. \ekvdefunknownNoVal and \ekvredirectunknownNoVal are mutually exclusive, you can't use both.

Example: See above.

1.2 Parsing Keys

\ekvset

```
\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath}\amb}\amb}\amb}}}}}}}}}}}}}}
```

Splits $\langle key \rangle = \langle value \rangle$ pairs on commas. From both $\langle key \rangle$ and $\langle value \rangle$ up to one space is stripped from both ends, if then only a braced group remains the braces are stripped as well. So $\langle bar = baz \rangle$ and $\langle bar = baz \rangle$ will both do $\langle bar = baz \rangle$, so you can hide commas, equal signs and spaces at the ends of either $\langle bar \rangle$ or $\langle bar \rangle$ by putting braces around them. If you omit the equal sign the code of the key created with the NoVal variants described in subsection 1.1 will be executed. If $\langle bar \rangle = \langle bar \rangle$ contains more than a single unhidden equal sign, it will be split at the first one and the others are considered part of the value. $\langle bar \rangle$ should be nestable.

\ekvset is currently *not* alignment safe.² As a result, key names and values that contain an & must be wrapped in braces when \ekvset is used inside an alignment (like LATEX 2ε 's tabular environment) or you have to create a wrapper that ensures an alignment safe context.

Example: Parse key=arg, key in the set foo: \ekvset { foo } { key=arg, key }

\ekvsetSneaked

```
\ensuremath{\verb| kvsetSneaked{\langle set \rangle} {\langle sneak \rangle} {\langle key \rangle} = \langle value \rangle, \ldots}
```

Just like \ekvset, this macro parses the $\langle key \rangle = \langle value \rangle$ pairs within the given $\langle set \rangle$. But \ekvsetSneaked will behave as if \ekvsneak has been called with $\langle sneak \rangle$ as its argument as the first action.

Example: Parse key=arg, key in the set foo with $\arrange key=arg$, key \arr

\ekvsetdef

```
\ensuremath{\mbox{ekvsetdef}\langle cs\rangle}{\langle set\rangle}
```

With this function you can define a shorthand macro $\langle cs \rangle$ to parse keys of a specified $\langle set \rangle$. It is always defined $\lceil long \rceil$, but if you need to you can also prefix it with $\lceil long \rceil$. The resulting macro is faster than but else equivalent to the idiomatic definition:

```
\label{eq:cs} $$  \log \left( cs \right) $$ 1{\left( set \right)} $$ $$ $$
```

Example: Define the macro \foosetup to parse keys in the set foo and use it to parse key=arg, key:

```
\ekvsetdef\foosetup{foo}
\foosetup{key=arg, key}
```

\ekvsetSneakeddef

```
\verb|\ekvsetSneakeddef| \langle cs \rangle \{ \langle set \rangle \}|
```

Just like \ekvsetdef this defines a shorthand macro $\langle cs \rangle$, but this macro will make it a shorthand for \ekvsetSneaked, meaning that $\langle cs \rangle$ will take two arguments, the first being stuff that should be given to \ekvsneak and the second the $\langle key \rangle = \langle value \rangle$ list. The resulting macro is faster than but else equivalent to the idiomatic definition:

Example: Define the macro \foothings to parse keys in the set foo and accept a sneaked argument, then use it to parse key=arg, key and sneak \afterwards:

²This might change in the future, I've not decided yet.

```
\ekvsetSneakeddef\foothings{foo}
\foothings{\afterwards}{key=arg, key}
```

\ekvsetdefSneaked

 $\verb|\ekvsetdefSneaked| \langle cs \rangle {\langle set \rangle} {\langle sneaked \rangle}|$

And this one behaves like \ekvsetSneakeddef but with a fixed \(sneaked \) argument. So the resulting macro is faster than but else equivalent to the idiomatic definition:

```
\label{cs} $$ \prod \ext{Sneaked}(\ent)}{\space{2mm} } $$
```

Example: Define the macro \barthing to parse keys in the set bar and always execute \afterwards afterwards, then use it to parse key=arg, key:

```
\ekvsetdefSneaked\barthing{bar}{\afterwards}\barthing{key=arg, key}
```

\ekvparse

```
\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath}\amb}\amb}\amb}}}}}}}}}}}}}}
```

This macro parses the $\langle key \rangle = \langle value \rangle$ pairs and provides those list elements which are only keys as an argument to $\langle code1 \rangle$, and those which are a $\langle key \rangle = \langle value \rangle$ pair to $\langle code2 \rangle$ as two arguments. It is fully expandable as well and returns each element of the parsed list in \unexpanded, which has no effect outside of an \expanded or \edge context. Also \ekvparse expands in exactly two steps of expansion. You can use multiple tokens in $\langle code1 \rangle$ and $\langle code2 \rangle$ or just a single control sequence name. In both cases the found $\langle key \rangle$ and $\langle value \rangle$ are provided as a brace group following them.

\ekvparse is alignment safe, meaning that you don't have to take any precautions if it is used inside an alignment context (like \LaTeX 2 tabular environment) and any key or value can contain an &.

\ekvbreak, \ekvsneak, and \ekvchangeset and their relatives don't work in \ekvparse. It is analogue to expl3's \keyval_parse:NNn, but not with the same parsing rules — \keyval_parse:NNn throws an error on multiple equal signs per $\langle key \rangle = \langle value \rangle$ pair and on empty $\langle key \rangle$ names in a $\langle key \rangle = \langle value \rangle$ pair, both of which \ekvparse doesn't deal with.

Example:

```
\ensuremath{\current \current \curren
```

```
\handlekeyval{S}{foo}{bar}\handlekey{S}{key}\handlekeyval{S}{baz}{zzz}
```

and afterwards \handlekey and \handlekeyval would have to further handle the \(\lambda key \rangle \). There are no macros like these two contained in \(\text{exp_v} \), you have to set them up yourself if you want to use \ekvparse (of course the names might differ). If you need the results of \ekvparse as the argument for another macro, you should use \expanded, or expand \ekvparse twice, as only then the input stream will contain the output above:

```
\expandafter\parse\expanded {\ekvparse} \k\kv{foo = bar, key, baz={zzz}} \} or $$ \exp andafter\expandafter \\ \exp andafter\expandafter \\ parse\ekvparse\k\kv{foo = bar, key, baz={zzz}} \} $$ would both expand to $$ parse\kv{foo}{bar}\k{key}\kv{baz}{zzz} $$
```

1.3 Other Macros

exply provides some other macros which might be of interest.

\ekvVersion \ekvDate

These two macros store the version and date of the package.

\ekvifdefined \ekvifdefinedNoVal

```
\label{eq:continuity} $$\left(\frac{\langle set}{\langle set}\right)}{\langle true}\left(\frac{\langle set}{\langle set}\right)}{\langle true}\right)} $$ \left(\frac{\langle set}{\langle set}\right)}{\langle true}\left(\frac{\langle set}{\langle set}\right)}{\langle true}\right)} $$
```

These two macros test whether there is a $\langle key \rangle$ in $\langle set \rangle$. It is false if either a hash table entry doesn't exist for that key or its meaning is $\$ relax.

Example: Check whether the key special is already defined in set foo, if it isn't input a file that contains more key definitions:

```
\ekvifdefined { foo } { special } { } { \input { foo . morekeys . tex } }
```

\ekvifdefinedset

 $\verb|\ekvifdefinedset{\langle set \rangle}{\langle true \rangle}{\langle false \rangle}|$

This macro tests whether $\langle set \rangle$ is defined (which it is if at least one key was defined for it). If it is $\langle true \rangle$ will be run, else $\langle false \rangle$.

Example: Check whether the set VeRyUnLiKeLy is already defined, if so throw an error, else do nothing:

```
\ekvifdefinedset{VeRyUnLiKeLy}
{\errmessage{VeRyUnLiKeLy already defined}}{}
```

\ekvbreak \ekvbreakPreSneak \ekvbreakPostSneak $\verb|\ekvbreak|{|}\langle after|\rangle|$

Gobbles the remainder of the current \ekvset macro and its argument list and reinserts \(\after \). So this can be used to break out of \ekvset. The first variant will also gobble anything that has been sneaked out using \ekvsneak or \ekvsneakPre, while \ekvbreakPreSneak will put \(\after \) before anything that has been smuggled and \ekvbreakPostSneak will put \(\after \) after the stuff that has been sneaked out.

Example: Define a key abort that will stop key parsing inside the set foo and execute \foo@aborted, or if it got a value \foo@aborted@with:

```
\ekvdefNoVal{foo}{abort}{\ekvbreak{\foo@aborted}}
\ekvdef{foo}{abort}{\ekvbreak{\foo@aborted@with{#1}}}
```

∖ekvsneak ∖ekvsneakPre $\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath}\amb}\amb}\amb}\amb}}}}}}}}}$

Puts (after) after the effects of \ekvset. The first variant will put (after) after any other tokens which might have been sneaked before, while \ekvsneakPre will put (after) before other smuggled stuff. This reads and reinserts the remainder of the current \ekvset macro and its argument list to do its job. After \ekvset has parsed the entire (key)=(value) list everything that has been \ekvsneaked will be left in the input stream. A small usage example is shown in subsubsection 1.4.3.

Example: Define a key secret in the set foo that will sneak out foo@secretly@sneaked: $\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{$

\ekvchangeset

 $\ensuremath{\mbox{ekvchangeset}}\ensuremath{\mbox{(new-set)}}$

Replaces the current set with $\langle new-set \rangle$, so for the rest of the current $\langle ekvset call$, that call behaves as if it was called with $\langle ekvset \langle new-set \rangle$. It is comparable to using $\langle key \rangle$. cd in pgfkeys.

Example: Define a key cd in set foo that will change to another set as specified in the value, if the set is undefined it'll stop the parsing and throw an error as defined in the macro \foo@cd@error:

```
\ekvdef{foo}{cd} {\ekvifdefinedset{#1}}{\ekvchangeset{#1}}{\ekvbreak{\foo@cd@error}}}
```

\ekvoptarg

 $\verb|\ekvoptarg{$\langle next\rangle$} {\langle default\rangle}|$

This macro will check for a following optional argument in brackets ([]) expandably. After the optional argument there has to be a mandatory one. The code in $\langle next \rangle$ should expect two arguments (the processed optional argument and the mandatory one). If there was an optional argument the result will be $\langle next \rangle \{\langle optional \rangle\} \langle mandatory \rangle$ (so the optional argument will be wrapped in braces, the mandatory argument will be untouched). If there was no optional argument the result will be $\langle next \rangle \{\langle default \rangle\} \{\langle mandatory \rangle\}$ (so the default will be used and the mandatory argument will be wrapped in braces after being read once – if it was already wrapped it is effectively unchanged).

\ekvoptarg expands in exactly two steps, grabs all the arguments only at the second expansion step, and is alignment safe. It has its limitations however. It can't tell the difference between [and {[}, so it doesn't work if the mandatory argument is a single bracket. Also if the optional argument should contain a nested closing bracket, the optional argument has to use nested braces like so: [{arg]ument}].

Example: Say we have a macro that should take an optional argument defaulting to 1: $\mbox{$\mbox{$newcommand$} \ oo {\mbox{$\mbox{$ekvoptarg$} \end{$}}}$

```
\newcommand\@foo[2]{Mandatory: #2\par Optional: #1}
```

\ekvoptargTF

 $\verb|\ekvoptargTF{| \langle true \rangle} { \langle false \rangle}|$

This macro is similar to \ekvoptarg, but will result in $\langle true \rangle \{\langle optional \rangle\} \langle mandatory \rangle$ or $\langle false \rangle \{\langle mandatory \rangle\}$ instead of placing a default value.

\ekvoptargTF expands in exactly two steps, grabs all the arguments only at the second expansion step, and is alignment safe. It has the same limitations as \ekvoptarg. *Example:* Say we have a macro that should behave differently depending on whether there was an optional argument or not. This could be done with:

```
\newcommand\foo@a\foo@b\\\newcommand\foo@a[2]{Mandatory: #2\par Optional: #1\\\newcommand\foo@b[1]{Mandatory: #1\par No optional.}
```

\ekvcsvloop

 $\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath{\ensuremath}\ensuremath{\ensuremath{\ens$

This loops over the comma separated items in $\langle csv-list \rangle$ and, after stripping spaces from either end of $\langle item \rangle$ and removing at most one set of outer braces, leaves $\mbox{unexpanded}\{\langle code \rangle\{\langle item \rangle\}\}$ for each list item in the input stream. Blank elements are ignored (if you need a blank element it should be given as $\{\}$). It supports both active commas and commas of category other. You could consider it as a watered down version of $\mbox{kevparse}$. However it is not alignment safe, which you could achieve by nesting it in $\mbox{kevparse}$ (since the braces around the argument of $\mbox{kevparse}$ will hide &s from \mbox{TeX} 's alignment parsing).

Example: The following splits a comma separated list and prints it in a typewriter font with parentheses around each element.

```
\newcommand*\myprocessor[1]{\texttt((#1))}
\ekvcsvloop\myprocessor{abc, def, ghi}\par
\ekvcsvloop\myprocessor{1,,2,,3,,4}\par
```

(abc)(def)(ghi) (1)(2)(3)(4)

\ekverr

```
\ensuremath{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath}\ensuremath{\ensuremath{\mbox{\ensuremath{\mbox{\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\engenturemath}\engenturemath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ensuremath}\ens
```

This macro will throw an error fully expandably.³ The error length is limited to a total length of 69 characters, and since ten characters will be added for the formatting (! $_{\square}$ and $_{\square}$ Error: $_{\square}$) that leaves us with a total length for $\langle package \rangle$ plus $\langle message \rangle$ of 59 characters. If the message gets longer T_EX will only display the first 69 characters and append \ETC. to the end.

Neither $\langle package \rangle$ nor $\langle message \rangle$ expand any further. Also $\langle package \rangle$ must not contain an explicit $\langle message \rangle$. No such restriction applies to $\langle message \rangle$.

If J is set up as the <text> which is the case in $\Delta_E X \ge$ but not in plain $T_E X$ by default) you can use that to introduce line breaks in your error message. However that doesn't change the message length limit.

After your own error message some further text will be placed. The formatting of that text will look good if ^J is the \newlinechar, else not so much. That text will read:

! Paragraph ended before \<an-expandable-macro> completed due to above exception. If the error summary is not comprehensible see the package documentation.

```
I will try to recover now. If you're in interactive mode hit <return> at the ? prompt and I continue hoping recovery was complete.
```

Any clean up has to be done by you, \ekverr will expand to nothing after throwing the error message.

In ConTEXt this macro works differently. While still being fully expandable, it doesn't have the character count limitation and doesn't impose restrictions on \(\lambda package \). It will not display the additional text and adding line breaks is not possible.

Example: Say we set up a small calculation which works with user input. In our calculation we need a division, so have to watch out for division by zero. If we detect such a case we throw an error and do the recovery by using the biggest integer allowed in TeX as the result.

```
\newcommand*\mydivision[2]
{%
\the\numexpr
\ifnum\numexpr#2=0 % space here on purpose
\ekverr{my}{division by 0. Setting result to 2147483647.}%
\2147483647%
\else
\((\frac{#1}{#2})\)%
\fi
```

 $^{^3}$ The used mechanism was to the best of my knowledge first implemented by Jean-François Burnol.

```
\relax
(10+5)/(3-3) approx mydivision \{10+5\}\{3-3\}$
If that code gets executed the following will be the terminal output
Runaway argument?
! my Error: division by 0. Setting result to 2147483647.
! Paragraph ended before \<an-expandable-macro>
completed due to above exception. If the error
summary is not comprehensible see the package
documentation.
I will try to recover now. If you're in inter-
active mode hit <return> at the ? prompt and I
continue hoping recovery was complete.
<to be read again>
                    \par
1.15 \ (10+5)/(3-3) \approx\mydivision{10+5}{3-3}
?
                            (10+5)/(3-3) \approx 2147483647
and the output would contain
                                                     if we continued the T<sub>E</sub>X run
at the prompt.
```

\ekv@name \ekv@name@set \ekv@name@key

```
\label{eq:constraint} $\ekv@name{\langle set \rangle} {\langle key \rangle} $$ \ekv@name@set{\langle set \rangle} $$ \ekv@name@key{\langle key \rangle} $$
```

The names of the macros that correspond to a key in a set are build with these macros. The name is built from two blocks, one that is formatting the $\langle set \rangle$ name (\ekv@name@set) and one for formatting the $\langle key \rangle$ name (\ekv@name@key). To get the actual name the argument to \ekv@name@key must be \detokenized. Both blocks are put together (with the necessary \detokenize) by \ekv@name. For NoVal keys an additional N gets appended irrespective of these macros' definition, so their name is \ekv@name{ $\langle set \rangle$ }{ $\langle key \rangle$ }N.

You can use these macros to implement additional functionality or access key macros outside of expkv, but don't change them! expkv relies on their exact definitions internally. Example: Execute the callback of the NoVal key key in set foo: \csname\ekv@name \{ foo \} \ key \}N\endcsname

1.4 Examples

1.4.1 Standard Use-Case

Say we have a macro for which we want to create a $\langle key \rangle = \langle value \rangle$ interface. The macro has a parameter, which is stored in the dimension \ourdim having a default value from its initialisation. Now we want to be able to change that dimension with the width key to some specified value. For that we'd do

```
\newdimen\ourdim
\ourdim=150pt
\protected\ekvdef{our}{width}{\ourdim=#1\relax}
```

as you can see, we use the set our here. We want the key to behave different if no value is specified. In that case the key should not use its initial value, but be smart and determine the available space from \hsize, so we also define

```
\protected \earline{\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\color=0.9\col
```

Now we set up our macro to use this $\langle key \rangle = \langle value \rangle$ interface

```
\protected\def\ourmacro#1%
{\begingroup\ekvset{our}{#1}\the\ourdim\endgroup}
```

Finally we can use our macro like in the following

```
\ourmacro{}\par
\ourmacro{width}\par
\ourmacro{width=5pt}\par
```

The same keys using expky | U

1.4.2 A Macro to Draw Rules

Another small example could be a $\langle key \rangle = \langle value \rangle$ driven \rule alternative, because I keep forgetting the correct order of its arguments. First we define the keys (and initialize the macros used to store the keys):

```
\makeatletter
\newcommand*\myrule@ht{1ex}
\newcommand*\myrule@wd{0.1em}
\newcommand*\myrule@raise{\z@}
\protected\ekvdef{myrule}{ht}{\def\myrule@ht{#1}}
\protected\ekvdef{myrule}{wd}{\def\myrule@wd{#1}}
\protected\ekvdef{myrule}{raise}{\def\myrule@raise{#1}}
\protected\ekvdef{myrule}{lower}{\def\myrule@raise{-#1}}
```

Then we define a macro to change the defaults outside of \myrule and \myrule itself:

```
\ekvsetdef\myruleset{myrule}
\newcommand*\myrule[1][]
{\begingroup\myruleset{#1}\myrule@out\endgroup}
```

And finally the output:

\newcommand*\myrule@out{\rule[\myrule@raise]\myrule@wd\myrule@ht}\makeatother

And we can use it:

```
a\myrule\par
a\myrule[ht=2ex,lower=.5ex]\par
\myruleset{wd=5pt}
a\myrule
```

```
aı
a|
a∎
```

1.4.3 An Expandable (key)=(value) Macro Using \ekvsneak

Let's set up an expandable macro, that uses a $\langle key \rangle = \langle value \rangle$ interface. The problems we'll face for this are:

- 1. ignoring duplicate keys
- 2. default values for keys which weren't used
- 3. providing the values as the correct argument to a macro (ordered)

First we need to decide which $\langle key \rangle = \langle value \rangle$ parsing macro we want to do this with, \ekvset or \ekvparse. For this example we also want to show the usage of \ekvsneak, hence we'll choose \ekvset. And we'll have to use \ekvset such that it builds a parsable list for our macro internals. To gain back control after \ekvset is done we have to put an internal of our macro at the start of that list, so we use an internal key that uses \ekvsneakPre after any user input.

To ignore duplicates will be easy if the value of the key used last will be put first in the list, so the following will use \ekvsneakPre for the user-level keys. If we wanted some key for which the first usage should be the binding one we would use \ekvsneak instead for that key.

Providing default values can be done in different ways, we'll use a simple approach in which we'll just put the outcome of our keys if they were used with default values before the parsing list terminator.

Ordering the keys can be done simply by searching for a specific token for each argument which acts like a flag, so our sneaked out values will include specific tokens acting as markers.

Now that we have answers for our technical problems, we have to decide what our example macro should do. How about we define a macro that calculates the sine of a number and rounds that to a specified precision? As a small extra this macro should understand input in radian and degree and the used trigonometric function should be selectable as well. For the hard part of this task (expandably evaluating trigonometric functions) we'll use the xfp package.

First we set up our keys according to our earlier considerations and set up the user facing macro \sine. The end marker of the parsing list will be a \sine@stop token, which we don't need to define and we put our defaults right before it. The user macro \sine uses \ekvoptargTF to check for the optional argument short cutting to the final step if no optional argument was found. This way we safe some time in this case, though we have to specify the default values twice.

```
\RequirePackage{xfp}
\makeatletter
\ekvdef{expex}{f}{\ekvsneakPre{\f{#1}}}}
```

```
\ekvdef{expex}{round}{\ekvsneakPre{\rnd{#1}}}
\ekvdefNoVal{expex}{degree}{\ekvsneakPre{\deg{d}}}
\ekvdefNoVal{expex}{radian}{\ekvsneakPre{\deg{}}}
\ekvdefNoVal{expex}{radian}{\ekvsneakPre{\deg{}}}
\ekvdefNoVal{expex}{internal}{\ekvsneakPre{\sine@rnd}}
\newcommand*\sine{\ekvoptargTF\sine@args{\sine@final{sin}{d}{3}}}
\newcommand*\sine@args[2]
{\ekvset{expex}{#1,internal}\rnd{3}\deg{d}\f{sin}\sine@stop{#2}}
```

Now we need to define some internal macros to extract the value of each key's last usage (remember that this will be the group after the first special flag-token). For that we use one delimited macro per key.

After the macros \sine@rnd, \sine@deg, and \sine@f the macro \sine@final will see \sine@final $\{\langle f \rangle\}$ {\langle degree/radian} \} {\langle round} \} {\langle round}. Now \sine@final has to expandably deal with those arguments such that the \fpeval macro of xfp gets the correct input. Luckily this is pretty straight forward in this example. In \fpeval the trigonometric functions have names such as \sin or \cos and the degree taking variants \sind or \cos d. And since the degree key puts a d in #2 and the radian key leaves #2 empty all we have to do to get the correct function name is stick the two together.

```
Let's test our macro:
```

The same macro using expkvics Using expkvics we can set up something equivalent with a bit less code. The implementation chosen in expkvics is more efficient than the example above and way easier to code for the user.

```
\makeatletter
\newcommand*\sine{\ekvoptargTF\sine@a{\sine@b{sin}{d}{3}}}
\ekvcSplitAndForward\sine@a\sine@b
{
    f=sin,
        unit=d,
        round=3,
    }
\ekvcSecondaryKeys\sine@a
{
    nmeta degree={unit=d},
        nmeta radian={unit={}},
    }
\newcommand*\sine@b[4]{\fpeval{round(#1#2(#4),#3)}}
\makeatother
```

The resulting macro will behave just like the one previously defined, but will have an additional unit key, since in expkvics every argument must have a value taking key which defines it.

1.5 Error Messages

expkv should only send messages in case of errors, there are no warnings and no info messages. In this subsection those errors are listed.

1.5.1 Load Time

expkv.tex checks whether ε -TeX and the \expanded primitive are available. If it isn't, an error will be thrown using \errmessage:

```
! expkv Error: e-TeX and \expanded required.
```

1.5.2 Defining Keys

If you get any error from expky while you're trying to define a key, the definition will be aborted and gobbled.

If you try to define a key with an empty set name you'll get:

```
! expkv Error: empty set name not allowed.
```

Similarly, if you try to define a key with an empty key name:

```
! expkv Error: empty key name not allowed.
```

Both of these messages are done in a way that doesn't throw additional errors due to \global, \long, etc., not being used correctly if you prefixed one of the defining macros.

1.5.3 Using Keys

This subsubsection contains the errors thrown during \ekvset. The errors are thrown in an expandable manner using \ekverr. In the following messages $\langle key \rangle$ gets replaced with the problematic key's name, and $\langle set \rangle$ with the corresponding set. If any errors during $\langle key \rangle = \langle value \rangle$ handling are encountered, the entry in the comma separated list will be omitted after the error is thrown and the next $\langle key \rangle = \langle value \rangle$ pair will be parsed.

If you're using an undefined key you'll get:

```
Runaway argument?
! expkv Error: unknown key '(key)' in set '(set)'
```

If you're using a key for which only a normal version and no NoVal version is defined, but don't provide a value, you'll get:

```
Runaway argument?
! expkv Error: missing value for '\langle key\' in set '\langle set\'
```

If you're using a key for which only a NoVal version and no normal version is defined, but provide a value, you'll get:

```
Runaway argument?
! expkv Error: unwanted value for '(key)' in set '(set)'
```

If you're using an undefined key in a set for which \ekvredirectunknown was used, and the key isn't found in any of the other sets as well, you'll get:

```
Runaway argument?
! expkv Error: no key '\langle key \rangle in sets \langle \langle set 1 \rangle \rangle \langle set 2 \rangle \rangle \...
```

If you're using an undefined NoVal key in a set for which \ekvredirectunknownNoVal was used, and the key isn't found in any of the other sets as well, you'll get:

```
Runaway argument?
! expkv Error: no NoVal key '(key)' in sets {(set1)}{(set2)}...
```

If you're using a set for which you never executed one of the defining macros from subsection 1.1 you'll get a low level TeX error, as that isn't actively tested by the parser (and hence will lead to undefined behaviour and not be gracefully ignored). The error will look like

```
! Missing \endcsname inserted.
<to be read again>
\! expkv Error: Set '\(set\)' undefined.
```

1.6 Bugs

Just like keyval, expkv is bug free. But if you find bugshidden features⁴ you can tell me about them either via mail (see the first page) or directly on GitHub if you have an account there: https://github.com/Skillmon/tex_expkv

1.7 Comparisons

This subsection makes some basic comparison between expkv and other $\langle key \rangle = \langle value \rangle$ packages. The comparisons are really concise, regarding speed, feature range (without listing the features of each package), and bugs and misfeatures.

Comparisons of speed are done with a very simple test key and the help of the l3benchmark package. The key and its usage should be equivalent to

```
\protected\ekvdef\test\{\height\}\\def\myheight\{#1\}\\ekvsetdef\expkvtest\test\}\\expkvtest\{\height} = 6\}
```

and only the usage of the key, not its definition, is benchmarked. For the impatient, the essence of these comparisons regarding speed and buggy behaviour is contained in Table 1.

As far as I know expkv is the only fully expandable $\langle key \rangle = \langle value \rangle$ parser. I tried to compare expkv to every $\langle key \rangle = \langle value \rangle$ package listed on CTAN, however, one might notice that some of those are missing from this list. That's because I didn't get the others to work due to bugs, or because they just provide wrappers around other packages in this list.

In this subsection is no benchmark of \ekvparse and \keyval_parse:NNn contained, as most other packages don't provide equivalent features to my knowledge. \ekvparse is slightly faster than \ekvset, but keep in mind that it does less. The same is true for \keyval_parse:NNn compared to \keys_set:nn of expl3 (where the difference is much

⁴Thanks, David!

bigger). Comparing just the two, \ekvparse is a tad faster than \keyval_parse:NNn because of the two tests (for empty key names and only a single equal sign) which are omitted.

keyval is about 30 % to 40 % faster and has a comparable feature set (actually a bit smaller since expkv supports unknown-key handlers and redirection to other sets) just a slightly different way how it handles keys without values. That might be considered a drawback, as it limits the versatility, but also as an advantage, as it might reduce doubled code. Keep in mind that as soon as someone loads xkeyval the performance of keyval gets replaced by xkeyval's.

Also keyval has a bug, which unfortunately can't really be resolved without breaking backwards compatibility for *many* documents, namely it strips braces from the argument before stripping spaces if the argument isn't surrounded by spaces, also it might strip more than one set of braces. Hence all of the following are equivalent in their outcome, though the last two lines should result in something different than the first two:

```
\setkeys{foo}{bar=baz}
\setkeys{foo}{bar= {baz}}
\setkeys{foo}{bar={ baz}} % should be 'baz'
\setkeys{foo}{bar={{baz}}} % should be '{baz}'
```

xkeyval is roughly twenty times slower, but it provides more functionality, e.g., it has choice keys, boolean keys, and so on. It contains the same bug as keyval as it has to be compatible with it by design (it replaces keyval's frontend), but also adds even more cases in which braces are stripped that shouldn't be stripped, worsening the situation.

Itxkeys is no longer compatible with the LATEX kernel starting with the release 2020-10-01. It is over 380 times slower – which is funny, because it aims to be "[...] faster [...] than these earlier packages [referring to keyval and xkeyval]." It needs more time to parse zero keys than five of the packages in this comparison need to parse 100 keys. Since it aims to have a bigger feature set than xkeyval, it most definitely also has a bigger feature set than expl.v. Also, it can't parse \long input, so as soon as your values contain a \par, it'll throw errors. Furthermore, Itxkeys doesn't strip outer braces at all by design, which, imho, is a weird design choice. In addition Itxkeys loads catoptions which is known to introduce bugs (e.g., see https://tex.stackexchange.com/questions/461783). Because it is no longer compatible with the kernel, I stop benchmarking it (so the numbers listed here and in Table 1 regarding Itxkeys were last updated on 2020-10-05).

l3keys is around four and a half times slower, but has an, imho, great interface to define keys. It strips *all* outer spaces, even if somehow multiple spaces ended up on either end. It offers more features, but is pretty much bound to expl3 code. Whether that's a drawback is up to you.

pgfkeys is around 2.7 times slower for one key if one uses the $\langle path \rangle$.cd syntax and almost 20% slower if one uses \pgfqkeys, but has an *enormous* feature set. To get the best performance \pgfqkeys was used in the benchmark. This reduces the overhead for setting the base directory of the benchmark keys by about 43 ops (so both p_0 and T_0 would be about 43 ops bigger if \pgfkeys{\(path \)/.cd,\(keys \)} was used instead). It has

the same or a very similar bug keyval has. The brace bug (and also the category fragility) can be fixed by pgfkeyx, but this package was last updated in 2012 and it slows down \pgfkeys by factor 8. Also pgfkeyx is no longer compatible with versions of pgfkeys newer than 2020-05-25.

kvsetkeys with kvdefinekeys is about 4.4 times slower, but it works even if commas and equals have category codes different from 12 (just as some other packages in this list). Else the features of the keys are equal to those of keyval, the parser has more features, though.

options is 1.7 times slower for only a single value. It has a much bigger feature set. Unfortunately it also suffers from the premature unbracing bug keyval has.

simplekv is hard to compare because I don't speak French (so I don't understand the documentation). There was an update released on 2020-04-27 which greatly improved the package's performance and adds functionality so that it can be used more like most of the other $\langle key \rangle = \langle value \rangle$ packages. It has problems with stripping braces and spaces in a hard to predict manner just like keyval. Also, while it tries to be robust against category code changes of commas and equal signs, the used mechanism fails if the $\langle key \rangle = \langle value \rangle$ list already got tokenised. Regarding unknown keys it got a very interesting behaviour. It doesn't throw an error, but stores the $\langle value \rangle$ in a new entry accessible with \useKV. Also if you omit $\langle value \rangle$ it stores true for that $\langle key \rangle$. For up to three keys, expkv is a bit faster, for more keys simplekv takes the lead.

YAX is over twenty times slower. It has a pretty strange syntax for the TEX-world, imho, and again a direct equivalent is hard to define (don't understand me wrong, I don't say I don't like the syntax, it's just atypical). It has the premature unbracing bug, too. Also somehow loading YAX broke options for me. The tested definition was:

```
\usepackage{yax}
\defactiveparameter yax {\storevalue\myheight yax:height } % setup
\setparameterlist{yax}{ height = 6 } % benchmark
```

1.8 License

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This work may be distributed and/or modified under the conditions of the LATEX Project Public License (LPPL), either version 1.3c of this license or (at your option) any later version. The latest version of this license is in the file:

```
http://www.latex-project.org/lppl.txt
This work is "maintained" (as per LPPL maintenance status) by
Jonathan P. Spratte.
```

Table 1: Comparison of $\langle key \rangle = \langle value \rangle$ packages. The packages are ordered from fastest to slowest for one $\langle key \rangle = \langle value \rangle$ pair. Benchmarking was done using l3benchmark and the scripts in the Benchmarks folder of the git repository. The columns p_i are the polynomial coefficients of a linear fit to the run-time, p_0 can be interpreted as the overhead for initialisation and p_1 the cost per key. The T_0 column is the actual mean ops needed for an empty list argument, as the linear fit doesn't match that point well in general. The column "BB" lists whether the parsing is affected by some sort of brace bug, "CF" stands for category code fragile and lists whether the parsing breaks with active commas or equal signs.

Package	p_1	p_0	T_0	BB	CF	Date
keyval	13.7	1.5	7.3	yes	yes	2014-10-28
expkv	19.7	2.2	6.6	no	no	2020-10-10
simplekv	18.3	7.0	17.7	yes	yes	2020-04-27
pgfkeys	24.3	1.7	10.7	yes	yes	2020-09-05
options	23.6	15.6	20.8	yes	yes	2015-03-01
kvsetkeys	*	*	40.3	no	no	2019-12-15
l3keys	71.3	33.1	31.6	no	no	2020-09-24
xkeyval	253.6	202.2	168.3	yes	yes	2014-12-03
YAX	421.9	157.0	114.7	yes	yes	2010-01-22
ltxkeys	3400.1	4738.0	5368.0	no	no	2012-11-17

^{*}For kvsetkeys the linear model used for the other packages is a poor fit, kvsetkeys seems to have approximately quadratic run-time, the coefficients of the second degree polynomial fit are $p_2 = 8.2$, $p_1 = 44.9$, and $p_0 = 60.8$. Of course the other packages might not really have linear run-time, but at least from 1 to 20 keys the fits don't seem too bad. If one extrapolates the fits for 100 $\langle key \rangle = \langle value \rangle$ pairs one finds that most of them match pretty well, the exception being ltxkeys, which behaves quadratic as well with $p_2 = 23.5$, $p_1 = 2906.6$, and $p_0 = 6547.5$.

2 Implementation

2.1 The LATEX Package

First we set up the LATEX package. That one doesn't really do much except \inputting the generic code and identifying itself as a package.

```
1 \def\ekv@tmp
2 {%
3   \ProvidesFile{expkv.tex}%
4     [\ekvDate\space v\ekvVersion\space an expandable key=val implementation]%
5 }
6 \input{expkv.tex}
7 \ProvidesPackage{expkv}%
8   [\ekvDate\space v\ekvVersion\space an expandable key=val implementation]
```

2.2 The ConTEXt module

This is pretty straight forward, we just have to change the error throwing mechanism for ConTEXt (the approach taken for plain and LATEX breaks in ConTEXt, effectively breaking ConTEXt, dropping you in an interactive TEX session with almost no means of escape).

```
9 \writestatus{loading}{ConTeXt User Module / expkv}
10 \unprotect
11 \input expkv.tex
12 \long\def\ekv@err@collect#1\par#2%
13 {\directlua{tex.error[[\detokenize{#2} Error: #1]]}}
14 \writestatus{loading}
15 {ConTeXt User Module / expkv / Version \ekvVersion\space loaded}
16 \protect\endinput
```

2.3 The Generic Code

The rest of this implementation will be the generic code.

We make sure that it's only input once:

```
17 \expandafter\ifx\csname ekvVersion\endcsname\relax
18 \else
    \expandafter\endinput
    Check whether \varepsilon-T<sub>E</sub>X and \expanded are available – exp<sub>k</sub>v requires \varepsilon-T<sub>E</sub>X.
21 \begingroup
    \edef\ekvtmpa{\string\expanded}
    \edef\ekvtmpb{\meaning\expanded}
    \expandafter
25 \endgroup
26 \ifx\ekvtmpa\ekvtmpb
    \expandafter\let\csname ekv@expanded\endcsname\expanded
    \expandafter\let\csname ekv@unexpanded\endcsname\unexpanded
29 \else
    \begingroup
       \edef\ekvtmpa{\string\expanded}
       \edef\ekvtmpb{\meaning\normalexpanded}
       \expandafter
    \endgroup
```

\ekvVersion \ekvDate

We're on our first input, so lets store the version and date in a macro.

```
44 \def\ekvVersion{1.9a}
45 \def\ekvDate{2021-09-20}
```

(End definition for \ekvVersion and \ekvDate. These functions are documented on page 7.)

If the LATEX format is loaded we want to be a good file and report back who we are, for this the package will have defined \ekv@tmp to use \ProvidesFile, else this will expand to a \relax and do no harm.

46 \csname ekv@tmp\endcsname

Store the category code of @ to later be able to reset it and change it to 11 for now.

- 47 \expandafter\chardef\csname ekv@tmp\endcsname=\catcode'\@
- 48 \catcode'\@=11

\ekv@tmp might later be reused to gobble any prefixes which might be provided to \ekvdef and similar in case the names are invalid, we just temporarily use it here as means to store the current category code of @ to restore it at the end of the file, we never care for the actual definition of it.

\ekv@if@lastnamedcs

If the primitive \lastnamedcs is available, we can be a bit faster than without it. So we test for this and save the test's result in this macro.

(End definition for \ekv@if@lastnamedcs.)

\ekv@empty

Sometimes we have to introduce a token to prevent accidental brace stripping. This token would then need to be removed by \@gobble or similar. Instead we can use \ekv@empty which will just expand to nothing, that is faster than gobbling an argument.

60 \def\ekv@empty{}

(End definition for \ekv@empty.)

\@gobble
\@firstofone
\@firstoftwo
\@secondoftwo
\ekv@fi@gobble
\ekv@fi@firstofone
\ekv@fi@firstoftwo
\ekv@fi@secondoftwo
\ekv@gobble@mark
\ekv@gobbleto@stop
\ekv@gobble@from@mark@to@stop

Since branching tests are often more versatile than \if...\else...\fi constructs, we define helpers that are branching pretty fast. Also here are some other utility functions that just grab some tokens. The ones that are also contained in LATEX don't use the ekv prefix. Not all of the ones defined here are really needed by expkv but are provided because packages like expkvider or expkvider need them (and I don't want to define them in each package which might need them).

```
61 \long\def\@gobble#1{}
62 \long\def\@firstofone#1{#1}
63 \long\def\@firstoftwo#1#2{#1}
64 \long\def\@secondoftwo#1#2{#2}
65 \long\def\ekv@fi@gobble\fi\@firstofone#1{\fi}
66 \long\def\ekv@fi@firstofone\fi\@gobble#1{\fi#1}
67 \long\def\ekv@fi@firstoftwo\fi\@secondoftwo#1#2{\fi#1}
68 \long\def\ekv@fi@secondoftwo\fi\@firstoftwo#1#2{\fi#2}
69 \def\ekv@gobble@mark\ekv@mark{}
70 \long\def\ekv@gobbleto@stop#1\ekv@stop{}
71 \long\def\ekv@gobble@from@mark@to@stop\ekv@mark#1\ekv@stop{}
```

(End definition for \@gobble and others.)

As you can see \ekv@gobbleto@stop uses a special marker \ekv@stop. The package will use three such markers, the one you've seen already, \ekv@mark and \ekv@nil. Contrarily to how for instance expl3 does things, we don't define them, as we don't need them to have an actual meaning. This has the advantage that if they somehow get expanded – which should never happen if things work out – they'll throw an error directly.

\ekv@ifempty@
\ekv@ifempty@true
\ekv@ifempty@false
\ekv@ifempty@true@F
\ekv@ifempty@true@F@gobble
\ekv@ifempty@true@F@gobbletwo

We can test for a lot of things building on an if-empty test, so lets define a really fast one. Since some tests might have reversed logic (true if something is not empty) we also set up macros for the reversed branches.

(End definition for \ekv@ifempty and others.)

\ekv@ifblank \ekv@ifblank@ The obvious test that can be based on an if-empty is if-blank, meaning a test checking whether the argument is empty or consists only of spaces. Our version here will be tweaked a bit, as we want to check this, but with one leading \ekv@mark token that is to be ignored. The wrapper \ekv@ifblank will not be used by expkv for speed reasons but expkvlopt uses it.

```
86 \long\def\ekv@ifblank#1%
87 {%
```

```
\lambda \ekv@ifblank@#1\ekv@nil\ekv@ifempty@B\ekv@ifempty@true
\lambda \ekv@ifempty@A\ekv@ifempty@B\@secondoftwo
\rangle \rang
```

(End definition for \ekv@ifblank and \ekv@ifblank@.)

\ekv@ifdefined

We'll need to check whether something is defined quite frequently, so why not define a macro that does this. The following test is expandable and pretty fast. The version with \lastnamedcs is the fastest version to test for an undefined macro I know of (that considers both undefined macros and those with the meaning \relax).

```
\ekv@if@lastnamedcs
93
       \long\def\ekv@ifdefined#1{\ifcsname#1\endcsname\ekv@ifdef@\fi\@secondoftwo}
       \def\ekv@ifdef@\fi\@secondoftwo
         {%
           \expandafter\ifx\lastnamedcs\relax
             \ekv@fi@secondoftwo
           \fi
           \@firstoftwo
    }
       \long\def\ekv@ifdefined#1%
105
106
           \ifcsname#1\endcsname\ekv@ifdef@\fi\ekv@ifdef@false#1\endcsname\relax
             \ekv@fi@secondoftwo
           \fi
           \@firstoftwo
       \def\ekv@ifdef@\fi\ekv@ifdef@false{\fi\expandafter\ifx\csname}
       \long\def\ekv@ifdef@false
           #1\endcsname\relax\ekv@fi@secondoftwo\fi\@firstoftwo#2#3%
         {#3}
    }
```

(End definition for \ekv@ifdefined.)

\ekv@strip@a \ekv@strip@b \ekv@strip@c We borrow some ideas of expl3's l3tl to strip spaces from keys and values. This \ekv@strip also strips one level of outer braces after stripping spaces, so an input of {abc} becomes abc after stripping. It should be used with #1 prefixed by \ekv@mark. Also this implementation at most strips one space from both sides (which should be fine most of the time, since TeX reads consecutive spaces as a single one during tokenisation).

```
117 \def\ekv@strip#1%
118 {%
119 \long\def\ekv@strip##1%
120 {%
121 \ekv@strip@a
122 ##1\ekv@nil
123 \ekv@mark#1%
124 #1\ekv@nil
125 }%
126 \long\def\ekv@strip@a##1\ekv@mark#1{\ekv@strip@b##1\ekv@mark}%
```

```
127  }
128 \ekv@strip{ }
129 \long\def\ekv@strip@b#1 \ekv@nil{\ekv@strip@c#1\ekv@nil}
130 \long\def\ekv@strip@c\ekv@mark#1\ekv@nil\ekv@mark#2\ekv@nil#3{#1}}
```

(End definition for \ekv@strip and others.)

\ekv@exparg@
\ekv@expargtwice
\ekv@expargtwice@
\ekv@zero

To reduce some code doublets while gaining some speed (and also as convenience for other packages in the family), it is often useful to expand the first token in a definition once. Let's define a wrapper for this.

Also, to end a \romannumeral expansion, we want to use \z@, which is contained in both plain TeX and LATeX, but we use a private name for it to make it easier to spot and hence easier to manage.

```
 \begin{array}{ll} & \text{lot} \\ & \text{long} \\ & \text{long}
```

(End definition for \ekv@exparg and others.)

\ekvcsvloop

\ekv@csv@loop@active \ekv@csv@loop@active@end An \ekvcsvloop will just loop over a csv list in a simple manner. First we split at active commas (gives better performance this way), next we have to check whether we're at the end of the list (checking for \ekv@stop). If not we go on splitting at commas of category other.

```
136 \begingroup
137 \def\ekvcsvloop#1{%
138 \endgroup
139 \long\def\ekvcsvloop##1##2%
140 {\ekv@csv@loop@active{##1}\ekv@mark##2#1\ekv@stop#1}
```

This does the same as \ekv@csv@loop but for active commas.

```
141 \long\def\ekv@csv@loop@active##1##2#1%
142 {%
143     \ekv@gobble@from@mark@to@stop##2\ekv@csv@loop@active@end\ekv@stop
144     \ekv@csv@loop{##1}##2,\ekv@stop,%
145  }%
146 \long\def\ekv@csv@loop@active@end
147     \ekv@stop
148     \ekv@stop
148     \ekv@csv@loop##1\ekv@mark\ekv@stop,\ekv@stop,%
149     {}%
150 }
```

Do the definitions with the weird catcode.

```
\catcode'\,=13
\catcode,
```

 $(\textit{End definition for } \verb+|ekv@csv@loop@active+, and \verb+|ekv@csv@loop@active@end+. These functions are documented on page 8.)$

\ekv@csv@loop@do \ekv@csv@loop@end We use temporary macros and an \expandafter chain to preexpand \ekv@strip here. After splitting at other commas we check again for end the end of the sublist, check for blank elements which should be ignored, and else strip spaces and execute the user code (protecting it from further expanding with \unexpanded).

```
153 \def\ekv@csv@loop#1%
```

```
\long\def\ekv@csv@loop##1##2,%
156
          {%
            \ekv@gobble@from@mark@to@stop##2\ekv@csv@loop@end\ekv@stop
            \ekv@ifblank@##2\ekv@nil\ekv@ifempty@B\ekv@csv@loop@blank
158
              \ekv@ifempty@A\ekv@ifempty@B
159
            #1\ekv@csv@loop@do{##1}%
160
          }%
161
     }
   \verb|\expandafter\ekv@csv@loop\expandafter{\ekv@strip{#2}}|
   \label{longdef} $$ \operatorname{long(def}(ekv@csv@loop@do#1#2{\ekv@unexpanded{#2{#1}}\ekv@csv@loop{#2}\ekv@mark}) $$
   \def\ekv@csv@loop@end#1%
166
        \long\def\ekv@csv@loop@end
167
            \ekv@stop
168
            \ekv@ifblank@\ekv@mark\ekv@stop\ekv@nil\ekv@ifempty@B\ekv@csv@loop@blank
169
            \ekv@ifempty@A\ekv@ifempty@B
            #1\ekv@csv@loop@do##1%
          {\ekv@csv@loop@active{##1}\ekv@mark}%
     }
   \expandafter\ekv@csv@loop@end\expandafter{\ekv@strip{\ekv@mark\ekv@stop}}
   \long\expandafter\def\expandafter\ekv@csv@loop@blank
        \expandafter\ekv@ifempty@A\expandafter\ekv@ifempty@B
176
       \ekv@strip{\ekv@mark#1}\ekv@csv@loop@do#2%
     {\ekv@csv@loop{#2}\ekv@mark}
178
(End definition for \ekv@csv@loop, \ekv@csv@loop@do, and \ekv@csv@loop@end.)
The keys will all follow the same naming scheme, so we define it here.
```

\ekv@name \ekv@name@set \ekv@name@key

```
179 \def\ekv@name@set#1{ekv#1(}
  \long\def\ekv@name@key#1{#1)}
181 \edef\ekv@name
     ₹%
182
183
       \ekv@unexpanded\expandafter{\ekv@name@set{#1}}%
       \ekv@unexpanded\expandafter{\ekv@name@key{\detokenize{#2}}}%
    }
\long\ekv@exparg{\def\ekv@name#1#2}{\ekv@name}
```

(End definition for \ekv@name, \ekv@name@set, and \ekv@name@key. These functions are documented on page 10.)

\ekv@undefined@set

We can misuse the macro name we use to expandably store the set-name in a single token - since this increases performance drastically, especially for long set-names to throw a more meaningful error message in case a set isn't defined. The name of \ekv@undefined@set is a little bit misleading, as it is called in either case inside of \csname, but the result will be a control sequence with meaning \relax if the set is undefined, hence will break the \csname building the key-macro which will throw the error message.

```
187 \def\ekv@undefined@set#1{! expkv Error: Set '#1' undefined.}
(End definition for \ekv@undefined@set.)
```

\ekv@checkvalid

We place some restrictions on the allowed names, though, namely sets and keys are not allowed to be empty - blanks are fine (meaning set- or key-names consisting of spaces). The \def\ekv@tmp gobbles any TFX prefixes which would otherwise throw errors. This

will, however, break the package if an \outer has been gobbled this way. I consider that good, because keys shouldn't be defined \outer anyways.

```
188 \edef\ekv@checkvalid
                                                                 {%
                                                                      \ekv@unexpanded\expandafter{\ekv@ifempty{#1}}%
                                                                      \ekv@unexpanded
                                                    191
                                                                          {{%
                                                     192
                                                                                \def\ekv@tmp{}%
                                                    193
                                                                                \errmessage{expkv Error: empty set name not allowed}%
                                                    194
                                                     195
                                                     196
                                                                                \ekv@unexpanded\expandafter{\ekv@ifempty{#2}}%
                                                     197
                                                                                \ekv@unexpanded
                                                                                    {%
                                                     199
                                                                                          {%
                                                                                               \def\ekv@tmp{}%
                                                                                               \errmessage{expkv Error: empty key name not allowed}%
                                                                                          }%
                                                                                          \@secondoftwo
                                                                                    }%
                                                                          }%
                                                                      \ekv@unexpanded{\@gobble}%
                                                                }
                                                           \ekv@exparg{\protected\def\ekv@checkvalid#1#2}{\ekv@checkvalid}%
                                                    (End definition for \ekv@checkvalid.)
                                                    And provide user-level macros to test whether a key is defined.
                                                    \ekv@expargtwice{\long\def\ekvifdefined#1#2}%
                                                                 \verb|\colored| \ensuremath{\texttt{NoVal#1#2}}| \ensure
                                                                (End definition for \ekvifdefined and \ekvifdefinedNoVal. These functions are documented on page 7.)
                                                   Set up the key defining macros \ekvdef etc. We use temporary macros to set these up
                             \ekvdef
                                                   with a few expansions already done.
                 \ekvdefNoVal
                             \ekvlet
                                                   214 \def\ekvdef#1#2#3#4%
                 \ekvletNoVal 215
                        \ekvletkv 216
                                                                      \protected\long\def\ekvdef##1##2##3%
                                                                           {#1{\expandafter\def\csname#2\endcsname###1{##3}#3}}%
            \ekvletkvNoVal 217
                                                                      \protected\long\def\ekvdefNoVal##1##2##3%
            \ekvdefunknown 218
                                                                           {#1{\expandafter\def\csname#2N\endcsname{##3}#3}}%
                                                   219
\ekvdefunknownNoVal
                                                                      \protected\long\def\ekvlet##1##2##3%
                                                                           {#1{\expandafter\let\csname#2\endcsname##3#3}}%
                                                                      \protected\long\def\ekvletNoVal##1##2##3%
                                                                           {#1{\expandafter\let\csname#2N\endcsname##3#3}}%
                                                                      \ekv@expargtwice{\protected\long\def\ekvdefunknown##1##2}%
                                                                           {%
                                                                                \romannumeral
                                                                                \expandafter\ekv@exparg@\expandafter
                                                                                    {%
                                                                                          \expandafter\expandafter\expandafter
```

\ekvifdefined \ekvifdefinedNoVal

\def\expandafter\csname\ekv@name{##1}{}u\endcsname###1###2{##2}%

```
#3%
            }%
            {\ekv@zero\ekv@checkvalid{##1}.}%
      \ekv@expargtwice{\protected\long\def\ekvdefunknownNoVal##1##2}%
        {%
236
          \romannumeral
          \expandafter\ekv@exparg@\expandafter
238
            {%
              \expandafter\expandafter\expandafter
              #3%
            }%
            {\ekv@zero\ekv@checkvalid{##1}.}%
245
      \protected\long\def\ekvletkv##1##2##3##4%
246
        {%
247
          #1%
248
249
              \expandafter\let\csname#2\expandafter\endcsname
              \csname#4\endcsname
              #3%
            }%
      \protected\long\def\ekvletkvNoVal##1##2##3##4%
        {%
          #1%
258
              \expandafter\let\csname#2N\expandafter\endcsname
259
              \csname#4N\endcsname
              #3%
            }%
        }%
263
    }
264
  \edef\ekvdefNoVal
265
266
      {\ekv@unexpanded\expandafter{\ekv@checkvalid{#1}{#2}}}%
267
      {\ekv@unexpanded\expandafter{\ekv@name{#1}{#2}}}%
268
269
         \ekv@unexpanded{\expandafter\ekv@defsetmacro\csname}%
         \ekv@unexpanded\expandafter{\ekv@undefined@set{#1}\endcsname{#1}}%
      }%
      {\tt \{\ekv@unexpanded\expandafter{\ekv@name{#3}{#4}}}\%
    }
  \expandafter\ekvdef\ekvdefNoVal
```

(End definition for \ekvdef and others. These functions are documented on page 3.)

\ekvredirectunknown \ekvredirectunknownNoVal

\ekv@defredirectunknown \ekv@redirectunknown@aux \ekv@redirectunknownNoVal@aux The redirection macros prepare the unknown function by looping over the provided list of sets and leaving a $\ensuremath{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\mbox{\m\s\m\s\s$

```
276 \protected\def\ekvredirectunknown
277 {%
278 \ekv@defredirectunknown
```

```
\ekv@redirect@kv
         \ekv@err@redirect@kv@notfound
280
         {\long\ekvdefunknown}%
281
         \ekv@redirectunknown@aux
282
283
   \protected\def\ekvredirectunknownNoVal
284
285
       \ekv@defredirectunknown
286
         \ekv@redirect@k
         \ekv@err@redirect@k@notfound
         \ekvdefunknownNoVal
         \ekv@redirectunknownNoVal@aux
     }
291
   \protected\def\ekv@defredirectunknown#1#2#3#4#5#6%
     {%
293
       \begingroup
294
       \edef\ekv@tmp
295
296
           \ekvcsvloop#1{#6}%
297
           \ekv@unexpanded{#2}%
           {\ekvcsvloop{}{#5,#6}}%
         }%
       \ekv@expargtwice
         {\endgroup#3{#5}}%
         {\expandafter#4\ekv@tmp\ekv@stop}%
  \def\ekv@redirectunknown@aux#1{#1{##1}{##2}}
  \def\ekv@redirectunknownNoVal@aux#1{#1{##1}}
```

(End definition for \ekvredirectunknown and others. These functions are documented on page 4.)

\ekv@redirect@k
\ekv@redirect@k@a
\ekv@redirect@k@a
\ekv@redirect@k@c
\ekv@redirect@k@c
\ekv@redirect@kw
\ekv@redirect@kv@a
\ekv@redirect@kv@a
\ekv@redirect@kv@a
\ekv@redirect@kv@a

The redirect code works by some simple loop over all the sets, which we already preprocessed in \ekv@defredirectunknown. For some optimisation we blow this up a bit code wise, essentially, all this does is \ekvifdefined or \ekvifdefinedNoVal in each set, if there is a match gobble the remainder of the specified sets and execute the key macro, else go on with the next set (to which the \(\key \) and \(\value \) are forwarded).

First we set up some code which is different depending on \lastnamedcs being available or not. All this is stored in a temporary macro to have pre-expanded \ekv@name constellations ready.

```
307 \def\ekv@redirect@k#1#2#3#4%
     {%
308
       \ekv@if@lastnamedcs
309
         {%
           \def\ekv@redirect@k##1##2##3%
             {%
312
                \ifcsname#1\endcsname\ekv@redirect@k@a\fi
               ##3{##1}%
314
           \def\ekv@redirect@k@a\fi{\fi\expandafter\ekv@redirect@k@b\lastnamedcs}%
           \long\def\ekv@redirect@kv##1##2##3##4%
                \ifcsname#2\endcsname\ekv@redirect@kv@a\fi\@gobble{##1}%
               ##4{##1}{##2}%
             }%
321
```

```
\def\ekv@redirect@kv@a\fi\@gobble
             {\fi\expandafter\ekv@redirect@kv@b\lastnamedcs}%
         }
         {%
           \def\ekv@redirect@k##1##2##3%
               \ifcsname#1\endcsname\ekv@redirect@k@a\fi\ekv@redirect@k@a@
                 #1\endcsname
329
               ##3{##1}%
             }%
           \def\ekv@redirect@k@a@#3\endcsname{}%
           \def\ekv@redirect@k@a\fi\ekv@redirect@k@a@
             {\fi\expandafter\ekv@redirect@k@b\csname}%
334
           \long\def\ekv@redirect@kv##1##2##3##4%
             {%
336
               \ifcsname#2\endcsname\ekv@redirect@kv@a\fi\ekv@redirect@kv@a@
                 #2\endcsname{##1}%
338
               ##4{##1}{##2}%
           \long\def\ekv@redirect@kv@a@#4\endcsname##3{}%
           \def\ekv@redirect@kv@a\fi\ekv@redirect@kv@a@
             {\fi\expandafter\ekv@redirect@kv@b\csname}%
         }%
    }
345
```

The key name given to this loop will already be \detokenized by \ekvset, so we can safely remove the \detokenize here for some performance gain.

```
346 \def\ekv@redirect@kv#1\detokenize#2#3\ekv@stop{\ekv@unexpanded{#1#2#3}}
347 \dedf\ekv@redirect@kv
348 {%
349 {\expandafter\ekv@redirect@kv\ekv@name{#2}{#1}N\ekv@stop}%
350 {\expandafter\ekv@redirect@kv\ekv@name{#3}{#2}\ekv@stop}%
351 {\expandafter\ekv@redirect@kv\ekv@name{#1}{#2}N\ekv@stop}%
352 {\expandafter\ekv@redirect@kv\ekv@name{#1}{#2}\ekv@stop}%
353 }
```

Everything is ready to make the real definitions.

354 \expandafter\ekv@redirect@k\ekv@redirect@kv

The remaining macros here are independent on \lastnamedcs, starting from the @b we know that there is a hash table entry, and get the macro as a parameter. We still have to test whether the macro is \relax, depending on the result of that test we have to either remove the remainder of the current test, or the remainder of the set list and invoke the macro.

```
355 \def\ekv@redirect@k@b#1%
356 {\ifx\relax#1\ekv@redirect@k@c\fi\ekv@redirect@k@d#1}
357 \def\ekv@redirect@k@c\fi\ekv@redirect@k@d#1{\fi}
358 \def\ekv@redirect@k@d#1#2\ekv@stop{#1}
359 \def\ekv@redirect@kv@b#1%
360 {\ifx\relax#1\ekv@redirect@kv@c\fi\ekv@redirect@kv@d#1}
361 \long\def\ekv@redirect@kv@c\fi\ekv@redirect@kv@d#1#2{\fi}
362 \long\def\ekv@redirect@kv@d#1#2#3\ekv@stop{#1{#2}}
```

(End definition for \ekv@redirect@k and others.)

\ekv@defsetmacro

In order to enhance the speed the set name given to \ekvset will be turned into a control sequence pretty early, so we have to define that control sequence.

```
363 \edef\ekv@defsetmacro
364 {%
365 \ekv@unexpanded{\ifx#1\relax\edef#1##1}%
366 {%
367 \ekv@unexpanded\expandafter{\ekv@name@set{#2}}%
368 \ekv@unexpanded\expandafter{\ekv@name@key{##1}}%
369 }%
370 \ekv@unexpanded{\fi}%
371 }
372 \ekv@exparg{\protected\def\ekv@defsetmacro#1#2}{\ekv@defsetmacro}
(End definition for \ekv@defsetmacro.)
```

\ekvifdefinedset

(End definition for \ekvifdefinedset. This function is documented on page 7.)

\ekvset

Set up \ekvset, which should not be affected by active commas and equal signs. The equal signs are a bit harder to cope with and we'll do that later, but the active commas can be handled by just doing two comma-splitting loops one at actives one at others. That's why we define \ekvset here with a temporary meaning just to set up the things with two different category codes. #1 will be a $_{13}$ and #2 will be a $_{13}$.

```
375 \begingroup
376 \def\ekvset#1#2{%
377 \endgroup
378 \ekv@exparg{\long\def\ekvset##1##2}%
379 {%
380 \expandafter\expandafter\expandafter
381 \ekv@set\expandafter\csname\ekv@undefined@set{##1}\endcsname
382 \ekv@mark##2#1\ekv@stop#1{}%
383 }
```

(End definition for \ekvset. This function is documented on page 5.)

\ekv@se

\ekv@set will split the $\langle key \rangle = \langle value \rangle$ list at active commas. Then it has to check whether there were unprotected other commas and resplit there.

```
384 \long\def\ekv@set##1##2#1%
385 {%
```

Test whether we're at the end, if so invoke \ekv@endset,

\ekv@gobble@from@mark@to@stop##2\ekv@endset\ekv@stop else go on with other commas.

```
387 \ekv@set@other##1##2,\ekv@stop,%
388 }
(End definition for \ekv@set.)
```

\ekv@endset

\ekv@endset is a hungry little macro. It will eat everything that remains of \ekv@set and unbrace the sneaked stuff.

```
389 \long\def\ekv@endset
390 \ekv@stop\ekv@set@other##1\ekv@mark\ekv@stop,\ekv@stop,##2%
391 {##2}
```

\ekv@eq@other \ekv@eq@active Splitting at equal signs will be done in a way that checks whether there is an equal sign and splits at the same time. This gets quite messy and the code might look complicated, but this is pretty fast (faster than first checking for an equal sign and splitting if one is found). The splitting code will be adapted for \ekvset and \ekvparse to get the most speed, but some of these macros don't require such adaptions. \ekv@eq@other and \ekv@eq@other and \ekv@eq@active will split the argument at the first equal sign and insert the macro which comes after the first following \ekv@mark. This allows for fast branching based on TeX's argument grabbing rules and we don't have to split after the branching if the equal sign was there.

```
\label{longdef} $$ \order{1=\#2\ekv@mark\#3{\#3\#1\ekv@stop\ekv@mark\#2} \order{\ekv@eq@active\#1\#2\#2\ekv@mark\#3{\#3\#1\ekv@stop\ekv@mark\#2} \ekv@mark\#2} $$
```

(End definition for \ekv@eq@other and \ekv@eq@active.)

\ekv@set@other

The macro \ekv@set@other is guaranteed to get only single $\langle key \rangle = \langle value \rangle$ pairs.

```
_{\rm 394} \long\def\ekv@set@other##1##2,% _{\rm 395} {%
```

First we test whether we're done.

(End definition for \ekv@endset.)

 $\verb|\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{|}\ensuremath{$

If not we split at the equal sign of category other.

```
\ekv@eq@other##2\ekv@nil\ekv@mark\ekv@set@eq@other@a
=\ekv@mark\ekv@set@eq@active
```

And insert the set name for the next recursion step of \ekv@set@other.

```
399 ##1%
400 \ekv@mark
401 }
(End definition for \ekv@set@other.)
```

\ekv@set@eq@other@a \ekv@set@eq@other@b The first of these two macros runs the split-test for equal signs of category active. It will only be inserted if the $\langle key \rangle = \langle value \rangle$ pair contained at least one equal sign of category other and ##1 will contain everything up to that equal sign.

The second macro will have been called by \ekv@eq@active if no active equal sign was found. All it does is remove the excess tokens of that test and forward the \(\key \) = \(\value \) pair to \ekv@set@pair. Normally we would have to also gobble an additional \ekv@mark after \ekv@stop, but this mark is needed to delimit \ekv@set@pair's argument anyway, so we just leave it there.

```
407 \ekv@exparg
```

```
408 {%
409 \long\def\ekv@set@eq@other@b
410 ##1\ekv@nil\ekv@mark\ekv@set@eq@other@active\ekv@stop\ekv@mark
411 ##2\ekv@nil=\ekv@mark\ekv@set@eq@active
412 }%
413 {\ekv@strip{##1}{\expandafter\ekv@set@pair\detokenize}\ekv@mark##2\ekv@nil}
```

(End definition for \ekv@set@eq@other@a and \ekv@set@eq@other@b.)

\ekv@set@eq@other@active

\ekv@set@eq@other@active will be called if the \(\lambda ey \rangle = \lambda value \rangle\) pair was wrongly split on an equal sign of category other but has an earlier equal sign of category active. ##1 will be the contents up to the active equal sign and ##2 everything that remains until the first found other equal sign. It has to reinsert the equal sign and forward things to \ekv@set@pair.

(End definition for \ekv@set@eq@other@active.)

\ekv@set@eq@active \ekv@set@eq@active@ \ekv@set@eq@active will be called when there was no equal sign of category other in the $\langle key \rangle = \langle value \rangle$ pair. It removes the excess tokens of the prior test and split-checks for an active equal sign.

```
421 \long\def\ekv@set@eq@active
422 ##1\ekv@nil\ekv@mark\ekv@set@eq@other@a\ekv@stop\ekv@mark
423 {%
424 \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@set@eq@active@
425 #2\ekv@mark\ekv@set@noeq
426 }
```

If an active equal sign was found in $\ensuremath{\mbox{\mbox{$\sim$}}}$ derive we'll have to pass the now split $(\ensuremath{\mbox{$key$}}) = (\ensuremath{\mbox{$value$}})$ pair on to $\ensuremath{\mbox{$\sim$}}$ derive depair.

```
427 \ekv@exparg
428 {\long\def\ekv@set@eq@active@##1\ekv@stop##2\ekv@nil#2\ekv@mark\ekv@set@noeq}%
429 {\ekv@strip{##1}{\expandafter\ekv@set@pair\detokenize}\ekv@mark##2\ekv@nil}
```

 $(End\ definition\ for\ \verb+\ekv@set@eq@active+\ and\ \verb+\ekv@set@eq@active@.)$

\ekv@set@noeq

If no active equal sign was found by \ekv@set@eq@active there is no equal sign contained in the parsed list entry. In that case we have to check whether the entry is blank in order to ignore it (in which case we'll have to gobble the set-name which was put after these tests by \ekv@set@other). Else this is a NoVal key and the entry is passed on to \ekv@set@key.

```
430 \edef\ekv@set@noeq
431 {%
432 \ekv@unexpanded
433 {%
434 \ekv@ifblank@##1\ekv@nil\ekv@ifempty@B\ekv@set@was@blank
435 \ekv@ifempty@A\ekv@ifempty@B
436 }%
```

```
\ekv@unexpanded\expandafter
         {\ekv@strip{##1}{\expandafter\ekv@set@key\detokenize}\ekv@mark}%
438
    }
439
  \ekv@exparg
440
     {%
441
       \long\def\ekv@set@noeq
442
           ##1\ekv@nil\ekv@mark\ekv@set@eq@active@\ekv@stop\ekv@mark
443
     }%
444
     {\ekv@set@noeq}
   \expandafter\def\expandafter\ekv@set@was@blank
       \expandafter\ekv@ifempty@A\expandafter\ekv@ifempty@B
       \ekv@strip{\ekv@mark##1}##2\ekv@mark
     {\ekv@set@other}
(End definition for \ekv@set@noeq.)
```

\ekv@endset@other

All that's left for \ekv@set@other is the macro which breaks the recursion loop at the end. This is done by gobbling all the remaining tokens.

(End definition for \ekv@endset@other.)

\ekvbreak PreSneak \ekvbreakPostSneak

Provide macros that can completely stop the parsing of \ekvset, who knows what it'll be useful for.

```
455 \long\def\ekvbreak##1##2\ekv@stop#1##3{##1}
456 \long\def\ekvbreakPreSneak ##1##2\ekv@stop#1##3{##1##3}
457 \long\def\ekvbreakPostSneak##1##2\ekv@stop#1##3{##3##1}
```

 $(\textit{End definition for } \verb+\ekvbreak+PreSneak+, and \verb+\ekvbreak+PostSneak+. These functions are documented on page 7.)$

\ekvsneak \ekvsneak Pre One last thing we want to do for \ekvset is to provide macros that just smuggle stuff after \ekvset's effects.

```
\label{longdef} $$ \log\left(\frac{\#1\#\#2\ekv@stop\#1\#\#3\{\#\#2\ekv@stop\#1\{\#\#3\#\#1\}\}} \\ \log\left(\frac{\#1\#\#2\ekv@stop\#1\#\#2}{\#2\ekv@stop\#1\#\#3}\}\right) $$
```

(End definition for \ekvsneak and \ekvsneakPre. These functions are documented on page 7.)

\ekvparse

Additionally to the \ekvset macro we also want to provide an \ekvparse macro, that has the same scope as \keyval_parse: NNn from expl3. This is pretty analogue to the \ekvset implementation, we just put an \unexpanded here and there instead of other macros to stop the \expanded on our output. The \unexpanded\expanded\{\ldots\right)} ensures that the material is in an alignment safe group at all time, and that it doesn't expand any further in an \edef or \expanded context.

```
460 \long\def\ekvparse##1##2##3%
461 {%
462 \ekv@unexpanded\ekv@expanded
463 {{\ekv@parse{##1}{##2}\ekv@mark##3#1\ekv@stop#1}}%
464 }
```

(End definition for $\ensuremath{\verb{\sc Vparse}}$. This function is documented on page 6.)

```
\ekv@parse
                              465 \long\def\ekv@parse##1##2##3#1%
                                    {%
                                      \ekv@gobble@from@mark@to@stop##3\ekv@endparse\ekv@stop
                              467
                                      \ekv@parse@other{##1}{##2}##3,\ekv@stop,%
                              468
                                    }
                              469
                              (End definition for \ekv@parse.)
              \ekv@endparse
                              470 \long\def\ekv@endparse
                                      \ekv@stop\ekv@parse@other##1\ekv@mark\ekv@stop,\ekv@stop,%
                              (End definition for \ekv@endparse.)
           \ekv@parse@other
                              473 \long\def\ekv@parse@other##1##2##3,%
                              474
                                      \ekv@gobble@from@mark@to@stop##3\ekv@endparse@other\ekv@stop
                              475
                                      \ekv@eq@other##3\ekv@nil\ekv@mark\ekv@parse@eq@other@a
                                        =\ekv@mark\ekv@parse@eq@active
                                      {##1}{##2}%
                              478
                                      \ekv@mark
                              479
                              480
                              (End definition for \ekv@parse@other.)
     \ekv@parse@eq@other@a
     \ekv@parse@eq@other@b
                              481 \long\def\ekv@parse@eq@other@a##1\ekv@stop
                              482
                                      \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@parse@eq@other@active
                              483
                                        #2\ekv@mark\ekv@parse@eq@other@b
                              484
                                    }
                              485
                                 \ekv@exparg
                                    {%
                                      \long\def\ekv@parse@eq@other@b
                                          ##1\ekv@nil\ekv@mark\ekv@parse@eq@other@active\ekv@stop\ekv@mark
                               489
                                          ##2\ekv@nil=\ekv@mark\ekv@parse@eq@active
                                    }%
                               491
                                    {\ekv@strip{##1}\ekv@parse@pair##2\ekv@nil}
                              (End\ definition\ for\ \verb|\ekv@parse@eq@other@a|\ and\ \verb|\ekv@parse@eq@other@b|)
\ekv@parse@eq@other@active
                                 \ekv@exparg
                                    {%
                              494
                                      \long\def\ekv@parse@eq@other@active
                              495
                                           ##1\ekv@stop##2\ekv@nil#2\ekv@mark
                              496
                                           \ekv@parse@eq@other@b\ekv@mark##3=\ekv@mark\ekv@parse@eq@active
                              497
                                    }%
                                    {\ekv@strip{##1}\ekv@parse@pair##2=##3}
                              (End definition for \ekv@parse@eq@other@active.)
```

```
\ekv@parse@eq@active
\ekv@parse@eq@active@
                       500 \long\def\ekv@parse@eq@active
                              ##1\ekv@nil\ekv@mark\ekv@parse@eq@other@a\ekv@stop\ekv@mark
                       501
                            {%
                       502
                              \ekv@eq@active##1\ekv@nil\ekv@mark\ekv@parse@eq@active@
                       503
                                #2\ekv@mark\ekv@parse@noeq
                       504
                            }
                       506 \ekv@exparg
                            {\long\def\ekv@parse@eq@active@##1\ekv@stop##2#2\ekv@mark\ekv@parse@noeq}%
                            {\ekv@strip{##1}\ekv@parse@pair##2}
                       (End definition for \ekv@parse@eq@active and \ekv@parse@eq@active@.)
      \ekv@parse@noeq
                          \edef\ekv@parse@noeq
                            {%
                              \ekv@unexpanded
                       511
                                  \ekv@ifblank@##1\ekv@nil\ekv@ifempty@B\ekv@parse@was@blank
                                    \ekv@ifempty@A\ekv@ifempty@B
                                \ekv@unexpanded\expandafter{\ekv@strip{##1}\ekv@parse@key}%
                       516
                            }
                          \ekv@exparg
                       518
                            {%
                       519
                              \long\def\ekv@parse@noeq
                                  ##1\ekv@nil\ekv@mark\ekv@parse@eq@active@\ekv@stop\ekv@mark
                            }%
                            {\ekv@parse@noeq}
                          \expandafter\def\expandafter\ekv@parse@was@blank
                              \expandafter\ekv@ifempty@A\expandafter\ekv@ifempty@B
                              \ekv@strip{\ekv@mark##1}\ekv@parse@key
                       526
                            {\ekv@parse@other}
                       (End\ definition\ for\ \verb+\ekv@parse@noeq.)
  \ekv@endparse@other
                       528 \long\def\ekv@endparse@other
                              \ekv@stop
                              \ekv@eq@other\ekv@mark\ekv@stop\ekv@nil\ekv@mark\ekv@parse@eq@other@a
                              =\ekv@mark\ekv@parse@eq@active
                            {\ekv@parse}
                       (End definition for \ekv@endparse@other.)
      \ekv@parse@pair
     \ekv@parse@pair@
                       {\ekv@strip{##2}\ekv@parse@pair@{##1}}
                          \long\def\ekv@parse@pair@##1##2##3##4%
                       535
                       536
                              \ekv@unexpanded{##4{##2}{##1}}%
                               \ekv@parse@other{##3}{##4}%
                       538
                            }
                       539
```

(End definition for \ekv@parse@pair and \ekv@parse@pair@.)

\ekv@parse@key

\ekvsetSneaked

This macro can be defined just by expanding \ekvsneak once after expanding \ekvset. To expand everything as much as possible early on we use a temporary definition.

```
549 \edef\ekvsetSneaked
550 {%
551 \ekv@unexpanded{\ekvsneak{#2}}%
552 \ekv@unexpanded\expandafter{\ekvset{#1}{#3}}%
553 }
554 \ekv@expargtwice{\long\def\ekvsetSneaked#1#2#3}{\ekvsetSneaked}
```

(End definition for \ekvsetSneaked. This function is documented on page 5.)

\ekvchangeset

Provide a macro that is able to switch out the current (set) in \ekvset. This operation allows something similar to pgfkeys's (key)/.cd mechanism. However this operation can be more expensive than /.cd as we can't just redefine some token to reflect this, but have to switch out the set expandably, so this works similar to the \ekvsneak macros reading and reinserting things, but it only has to read and reinsert the remainder of the current key's replacement code.

```
555 \ekv@exparg{\def\ekvchangeset#1}%
556 {%
557 \expandafter\expandafter
558 \ekv@changeset\expandafter\csname\ekv@undefined@set{#1}\endcsname\ekv@empty
559 }
```

(End definition for \ekvchangeset. This function is documented on page 8.)

\ekv@changeset

This macro does the real change-out of \ekvchangeset. #2 will have a leading \ekv@empty so that braces aren't stripped accidentally, but that will not hurt and just expand to nothing in one step.

(End definition for \ekv@changeset.)

\ekv@set@pair@a \ekv@set@pair@b \ekv@set@pair@c \ekv@set@pair@d \ekv@set@pair@d \ekv@set@pair gets invoked with the space and brace stripped and \detokenized keyname as its first, the value as the second, and the set name as the third argument. It provides tests for the key-macros and everything to be able to throw meaningful error messages if it isn't defined. We have two routes here, one if \lastnamedcs is defined and one if it isn't. The big difference is that if it is we can omit a \csname and instead just expand \lastnamedcs once to get the control sequence. If the macro is defined the value

will be space and brace stripped and the key-macro called. Else branch into the error handling provided by \ekv@set@pair.

```
\ekv@if@lastnamedcs
     {%
562
       \long\def\ekv@set@pair#1\ekv@mark#2\ekv@nil#3%
563
564
           \ifcsname #3{#1}\endcsname\ekv@set@pair@a\fi\@secondoftwo
565
566
567
                \ifcsname #3{}u\endcsname\ekv@set@pair@a\fi\@secondoftwo
568
                  {#2}%
569
                  {%
                    \ekv@ifdefined{#3{#1}N}%
                      \ekv@err@noarg
                      \ekv@err@unknown
                      #3%
                  }%
                  {#1}%
           \ekv@set@other#3%
578
         }
579
       \def\ekv@set@pair@a\fi\@secondoftwo
580
         {\fi\expandafter\ekv@set@pair@b\lastnamedcs}
581
     }
582
583
       \long\def\ekv@set@pair#1\ekv@mark#2\ekv@nil#3%
584
585
            \ifcsname #3{#1}\endcsname
586
                \ekv@set@pair@a\fi\ekv@set@pair@c#3{#1}\endcsname
587
              {#2}%
588
589
                \ifcsname #3{}u\endcsname
590
                    \ekv@set@pair@a\fi\ekv@set@pair@c#3{}u\endcsname
591
                  {#2}%
592
                  {%
593
                    \ekv@ifdefined{#3{#1}N}%
                      \ekv@err@noarg
                      \ekv@err@unknown
                      #3%
                  }%
598
                  {#1}%
599
            \ekv@set@other#3%
601
       \def\ekv@set@pair@a\fi\ekv@set@pair@c{\fi\expandafter\ekv@set@pair@b\csname}
       \long\def\ekv@set@pair@c#1\endcsname#2#3{#3}
606 \long\def\ekv@set@pair@b#1%
     {%
607
       \int x#1\relax
608
         \ekv@set@pair@e
609
610
       \ekv@set@pair@d#1%
611
```

(End definition for \ekv@set@pair and others.)

\ekv@set@key@a \ekv@set@key@b \ekv@set@key@c Analogous to \ekv@set@pair, \ekv@set@key builds the NoVal key-macro and provides an error-branch. \ekv@set@key@ will test whether the key-macro is defined and if so call it, else the errors are thrown.

```
\ekv@if@lastnamedcs
616
       \long\def\ekv@set@key#1\ekv@mark#2%
617
618
            \ifcsname #2{#1}N\endcsname\ekv@set@key@a\fi\@firstofone
619
                \ifcsname #2{}uN\endcsname\ekv@set@key@a\fi\@firstofone
621
622
                  ₹%
                    \ekv@ifdefined{#2{#1}}%
623
                       \ekv@err@reqval
624
                       \ekv@err@unknown
625
                       #2%
626
                  }%
                  {#1}%
              }%
            \ekv@set@other#2%
630
631
       \def\ekv@set@key@a\fi\@firstofone{\fi\expandafter\ekv@set@key@b\lastnamedcs}
632
     }
633
     {%
634
       \long\def\ekv@set@key#1\ekv@mark#2%
635
636
            \ifcsname #2{#1}N\endcsname
637
                \ekv@set@key@a\fi\ekv@set@key@c#2{#1}N\endcsname
              {%
                \ifcsname #2{}uN\endcsname
                    \ekv@set@key@a\fi\ekv@set@key@c#2{}uN\endcsname
                  {%
642
                    \ekv@ifdefined{#2{#1}}%
643
                       \ekv@err@reqval
                       \ekv@err@unknown
645
                       #2%
646
                  }%
647
                  {#1}%
              }%
            \ekv@set@other#2%
       \def\ekv@set@key@a\fi\ekv@set@key@c{\fi\expandafter\ekv@set@key@b\csname}
652
       \long\def\ekv@set@key@c#1N\endcsname#2{#2}
653
     }
654
   \long\def\ekv@set@key@b#1%
655
     {%
656
       \ifx#1\relax
657
          \ekv@fi@secondoftwo
658
659
       \fi
       \@firstoftwo#1%
```

```
661 }
(End definition for \ekv@set@key and others.)
```

\ekvsetdef

Provide a macro to define a shorthand to use \ekvset on a specified \(\set \). To gain the maximum speed \ekvset is expanded twice by \ekv@exparg so that during runtime the macro storing the set name is already built and one \expandafter doesn't have to be used.

```
662 \ekv@expargtwice{\protected\def\ekvsetdef#1#2}%
663 {%
664 \romannumeral
665 \ekv@exparg{\ekv@zero\ekv@exparg{\long\def#1##1}}%
666 {\ekvset{#2}{##1}}%
667 }
```

(End definition for \ekvsetdef. This function is documented on page 5.)

\ekvsetSneakeddef \ekvsetdefSneaked

And do the same for \ekvsetSneaked in the two possible ways, with a fixed sneaked argument and with a flexible one.

```
668 \ekv@expargtwice{\protected\def\ekvsetSneakeddef#1#2}%
669
 670
                                                      \ensuremath{\verb||} \ensuremath{\ensuremath{||} \ensuremath{\ensuremat
 671
                                                                       {\ekvsetSneaked{#2}{##1}{##2}}%
 672
                                     }
                      \ekv@expargtwice{\protected\def\ekvsetdefSneaked#1#2#3}%
 675
                                                       \romannumeral
 676
                                                      \ekv@exparg{\ekv@zero\ekv@exparg{\long\def#1##1}}%
 677
                                                                       {\ekvsetSneaked{#2}{#3}{##1}}%
 678
 679
```

(End definition for \ekvsetSneakeddef and \ekvsetdefSneaked. These functions are documented on page 5.)

\ekv@alignsafe \ekv@endalignsafe

\ekvoptargTF

These macros protect the usage of ampersands inside of alignment contexts.

```
680 \begingroup
681 \catcode'\^^@=2
682 \Offirstofone{\endgroup
683 \def\ekv@alignsafe{\romannumeral\iffalse{\fi'^^@}}
684 }
685 \def\ekv@endalignsafe{\ifnum'{=\ekv@zero}\fi}
```

(End definition for \ekv@alignsafe and \ekv@endalignsafe.)

\ekvoptarg Provide macros to expan

Provide macros to expandably collect an optional argument in brackets. The macros here are pretty simple in nature compared to xparse's possibilities (they don't care for nested bracket levels).

We start with a temporary definition to pre-expand \ekv@alignsafe (will be #1) and \ekv@endalignsafe (will be #2). As \ekv@alignsafe starts with a \romannumeral we use that to also control the number of steps needed instead of adding another \romannumeral. For this we have to remove the space token from the end of \ekv@alignsafe's definition.

```
686 \begingroup
687 \def\ekvoptarg#1#2{%
688 \endgroup
```

The real definition starts an expansion context and afterwards grabs the arguments. #1 will be the next step, #2 the default value, and #3 might be an opening bracket, or the mandatory argument. We check for the opening bracket, if it is found grab the optional argument, else leave #1{#2} in the input stream after ending the expansion context.

```
689 \def\ekvoptarg{#1\ekv@optarg@a}
690 \long\def\ekv@optarg@a##1##2##3%
691 {%
692 \ekv@optarg@if\ekv@mark#3\ekv@mark\ekv@optarg@b\ekv@mark[\ekv@mark
693 #2%
694 \dfirstofone{ ##1}{##2}{##3}%
695 }%
```

The other variant of this will do roughly the same. Here, #1 will be the next step if an optional argument is found, #2 the next step else, and #3 might be the opening bracket or mandatory argument.

The two macros to grab the optional argument have to remove the remainder of the test and the wrong next step as well as grabbing the argument.

```
\long\def\ekv@optarg@b\ekv@mark[\ekv@mark\ifnum'##1\fi\@firstofone##2##3##4##5]%
\[
\{\text{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\mathrm{\matr\mathrm{\mathrm{\matr\m{\mathrm{\matr\m{\mathrm{\mathrm{\matr\m{\mathrm{\math
```

Do the definitions and add the test macro. We use \ekv@strip to remove the trailing space from the definition of \ekv@alignsafe.

```
700 \ekv@exparg
710 {%
711 \expandafter\ekv@strip\expandafter
712 {\expandafter\ekv@mark\ekv@alignsafe}%
713 \ekvoptarg
714 }%
715 \ekv@endalignsafe
716 \long\def\ekv@optarg@if#1\ekv@mark[\ekv@mark{}
```

(End definition for \ekvoptarg and \ekvoptargTF. These functions are documented on page 8.)

\ekverr

\ekv@err@collect \ekv@err@cleanup Since \ekvset is fully expandable as long as the code of the keys is (which is unlikely) we want to somehow throw expandable errors, in our case via a runaway argument (to my knowledge the first version of this method was implemented by Jean-François Burnol, many thanks to him). The first step is to ensure that the second argument (which might contain user input) doesn't contain tokens we use as delimiters (in this case \par), this will be done by the front facing macro \ekverr. But first we set some other things up.

We use a temporary definition for \ekverr to get multiple consecutive spaces. Then we set up the macro that will collect the error and the macro that will throw the error.

The latter will have an unreasonable long name. This way we can convey more information. Though the information in the macro name is static and has to be somewhat general to fit every occurence. The important bit is that the long named macro has a delimited argument and is short which will throw the error at the \par at the end of \ekv@err@collect. This macro has the drawback that it will only print nicely if the \newlinechar is ^^J.

717 \def\ekv@err@cleanup\par{}

```
718 \def\ekv@err@collect#1%
                 719
                        \def\ekv@err@collect##1\par##2%
                          {%
                            \expandafter
                            \ekv@err@cleanup
                            #1! ##2 Error: ##1\par
                        \def#1##1\thanks@jfbu{}%
                 726
                      }
                 728 \def\ekverr{ }
                 \expandafter\ekv@err@collect\csname <an-expandable-macro>^^J%
                      completed due to above exception. \ekverr If the error^^J%
                      summary is \ekverr not comprehensible \ekverr see the package^^J%
                      documentation. ^ J%
                      I will try to recover now. \ekverr If you're in inter-^^J%
                      active mode hit <return> \ekverr at the ? prompt and I^^J%
                      continue hoping recovery\endcsname
                 736 \long\def\ekverr#1#2{\expandafter\ekv@err@collect\detokenize{#2}\par{#1}}
                 (End definition for \ekverr, \ekv@err@collect, and \ekv@err@cleanup. These functions are documented on
                 We define a shorthand to throw errors in exply.
       \ekv@err
                 737 \ekv@exparg{\long\def\ekv@err#1}{\ekverr{expkv}{#1}}
                 (End definition for \ekv@err.)
                 Now we can use \ekv@err to set up some error messages so that we can later use those
\ekv@err@common
\ekv@err@common@
                 instead of the full strings.
\ekv@err@unknown
                 738 \long\def\ekv@err@common #1#2{\expandafter\ekv@err@common@\string#2{#1}}
 \ekv@err@noarg 739 \ekv@exparg{\long\def\ekv@err@common@#1'#2' #3.#4#5}%
                      {\ekv@err{#4 '#5' in set '#2'}}
\ekv@err@reqval
                 744 \ekv@exparg{\long\def\ekv@err@unknown#1}{\ekv@err@common{unknown key}{#1}}
                 742 \ekv@exparg{\long\def\ekv@err@noarg #1}
                      {\ekv@err@common{unwanted value for}{#1}}
                 745 \ekv@exparg{\long\def\ekv@err@redirect@kv@notfound#1#2#3\ekv@stop}%
                      {\ekv@err{no key '#2' in sets #3}}
                 747 \ekv@exparg{\def\ekv@err@redirect@k@notfound#1#2\ekv@stop}%
                      {\ekv@err{no NoVal key '#1' in sets #2}}
                 (End definition for \ekv@err@common and others.)
                      Now everything that's left is to reset the category code of @.
                 749 \catcode'\@=\ekv@tmp
```

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